F.3.4.4.6 Derived reset operator

REPLACE

— (**reject_on** (b) p) $\equiv$ (**not** (**accept_on** (b) **not** p)) .

— (**reject_on** (b) P) $\equiv$ (**not** (**accept_on** (b) **not** P)) .

WITH

— ~~(**reject_on** (b) p) $\equiv$ (**not** (**accept_on** (b) **not** p)) .~~

— (**reject_on** (b) P) $\equiv$ (**not** (**accept_on** (b) **not** P)) .

F.5.3.1 Neutral Satisfaction

REPLACE

— $w \models$ (**accept_on** (b) P) iff either

  — $w \models P$, or

  — For some $0 \leq i < |w|$, $w^i \models b$ and $w^{0,i-1}\top^\omega \models P$.
  Here, $w^{0,-1}$ denotes the empty word.

— $w \models$ (**accept_on** (b) P) iff either $w \models P$, or for some $0 \leq i < |w|$, $w^i \models b$ and $w^{0,i-1}\top^\omega \models P$.
A word satisfies property **accept_on** (b) P) if and only if P succeeds or if during the evaluation of P the expression b evaluates to true.

WITH

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A word satisfies property **accept_on** (b) P) if and only if P succeeds or if during the evaluation of P the expression b evaluates to true.~~

F.5.3.3 Vacuity

REPLACE

— $w \models^{\text{non}}$ **disable iff** $(b)P$ iff $w \models^{\text{non}} P$ and one of the following holds:

1. For every $0 \leq i < |w|, w^i \not\models b$.
2. For some prefix x of w, we have that for every $0 \leq i < |x|, x^i \not\models b$,and either $x\bot^\omega \models P$ or $x\top^\omega \not\models P$.

— $w \models^{\text{non}}$ **accept_on** (b) P iff $w \models^{\text{non}} P$ and one of the following holds:

1. For every $0 \leq i < |w|, w^i \not\models b$.
2. For some prefix x of w, we have that for every $0 \leq i < |x|, x^i \not\models b$,and either $x\bot^\omega \models P$ or $x\top^\omega \not\models P$.

— $w \models^{\text{non}}$ **reject_on** (b) P iff $w \models^{\text{non}} P$ and one of the following holds:

1. For every $0 \leq i < |w|, w^i \not\models b$.
2. For some prefix x of w, we have that for every $0 \leq i < |x|, x^i \not\models b$,and either $x\bot^\omega \models P$ or $x\top^\omega \not\models P$.

WITH

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— $w \models^{\text{non}}$ **accept_on** (b) P iff $w \models^{\text{non}} P$ and one of the following holds:

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1. For every $0 \leq i < |w|, w^i \not\models b$.
2. ~~For some prefix x of w, we have that~~ There exists a prefix x of w, such that for every $0 \leq i < |x|, x^i \not\models b$,and either $x\bot^\omega \models P$ or $x\top^\omega \not\models P$.

F.5.6.1 Neutral Satisfaction

REPLACE

- $w, L_0 \models$ (**accept_on** $(b)\, P$) iff either
  - $w, L_0 \models P$ and no letter of w satisfies b, or
  - For some $0 \leq i < |w|$, $w^i \models b$ and $w^{0,i-1}\top^\omega, L_0 \models P$. Here, $w^{0,-1}$ denotes the empty word.
- $w, L_0 \models$ (**accept_on** $(b)\, P$) iff either $w, L_0 \models P$ and no letter of w satisfies b, or for some $0 \leq i < |w|$, $w^i \models b$ and $w^{0,i-1}\top^\omega, L_0 \models P$.

WITH

- $w, L_0 \models$ (**accept_on** $(b)\, P$) iff either
  - $w, L_0 \models P$ and no letter of w satisfies b, or
  - For some $0 \leq i < |w|$, $w^i \models b$ and $w^{0,i-1}\top^\omega, L_0 \models P$. Here, $w^{0,-1}$ denotes the empty word.
- ~~$w, L_0 \models$ (**accept_on** $(b)\, P$) iff either $w, L_0 \models P$ and no letter of w satisfies b, or for some $0 \leq i < |w|$, $w^i \models b$ and $w^{0,i-1}\top^\omega, L_0 \models P$.~~